#### CO3091 - Computational Intelligence and Software Engineering

Lecture 12



#### Multi-Objective Evolutionary Algorithms

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#### Overview

- Multi-objective optimisation problems.
- Single-objective algorithmic design for multi-objective problems.
- Non-dominated Sorting Genetic Algorithm II (NSGA-II).

#### Multi-Objective Optimisation Problems

- Real world problems frequently have more than one objective.
- Software project scheduling problem:
  - Cost and duration (to be minimised).
- Some requirements selection formulations:
  - Cost (to be minimised) and score (to be maximised).

#### Using Single-Objective Algorithms for Multi-Objective Problems

- In order to deal with multi-objective problems by using singleobjective optimisation algorithms, the multiple objectives have to be combined into a single fitness function.
- E.g.: software project scheduling problem

 $fitness(\mathbf{x}) = W_{cost} * Cost(\mathbf{x}) + W_{dur} * duration(\mathbf{x})$ (to be minimised)

W<sub>cost</sub> and W<sub>dur</sub>  $\in [0,1]$ W<sub>cost</sub> + W<sub>dur</sub> = 1

#### Using Single-Objective Algorithms for Multi-Objective Problems

• Generic fitness function for *k* objectives:

$$fitness(x) = \sum_{i=1}^{k} W_i f_i(x)$$

Weighted average of objective functions f<sub>i</sub>(x).

 $W_i \in [0, 1], \ 1 \le i \le k$ 

$$\sum_{i=1}^{k} W_i = 1$$

- Advantage of combining multiple objectives into a single fitness function:
  - Once weights are set, any single-objective optimisation algorithm can be used.

- Disadvantage 1:
  - Finding suitable weights is not easy before knowing what different trade-offs among objectives may be available.
  - Example:

fitness( $\mathbf{x}$ ) = W<sub>cost</sub> \* cost( $\mathbf{x}$ ) + W<sub>dur</sub> \* duration( $\mathbf{x}$ ) (to be minimised) W<sub>cost</sub> = 0.7 W<sub>dur</sub> = 0.3

Candidate solution 1: cost = 15,000 duration = 12 months

Candidate solution 2: cost = 15,500 duration = 24 months

Which of these solutions is better?

- Disadvantage 1:
  - Finding suitable weights is not easy before knowing what different trade-offs among objectives may be available.
  - Example:

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Candidate solution 1: cost = 15,000 duration = 12 months

fitness = 10503.6

Candidate solution 2: cost = 15,500duration = 24 months

fitness = 10857.2

- Disadvantage 1:
  - Finding suitable weights is not easy before knowing what different trade-offs among objectives may be available.
  - Example:

fitness( $\mathbf{x}$ ) = W<sub>cost</sub> \* cost( $\mathbf{x}$ ) + W<sub>dur</sub> \* duration( $\mathbf{x}$ ) (to be minimised) W<sub>cost</sub> = 0.7 W<sub>dur</sub> = 0.3

Candidate solution 1: cost = 15,000 duration = 12 months

Candidate solution 2: cost = 14,500 duration = 24 months

Which of these solutions is better?

- Disadvantage 1:
  - Finding suitable weights is not easy before knowing what different trade-offs among objectives may be available.
  - Example:

fitness( $\mathbf{x}$ ) = W<sub>cost</sub> \* cost( $\mathbf{x}$ ) + W<sub>dur</sub> \* duration( $\mathbf{x}$ ) (to be minimised) W<sub>cost</sub> = 0.7 W<sub>dur</sub> = 0.3

Candidate solution 1: cost = 15,000 duration = 12 months

fitness = 10503.6

Candidate solution 2: cost = 14,500 duration = 24 months

```
fitness = 10157.2
```

#### • Disadvantage 2:

- If the different objectives have different scales, weights must reflect not only preferences towards certain objectives, but also treat the different scales.
- E.g.: costs may typically be between 1000 and 80,000. Durations may typically vary between 3 and 48.

If  $W_{cost} = W_{dur} = 0.5$ , costs will tend to have greater importance.

- Disadvantage 3:
  - The weights will strongly affect the results.
    - You may not get certain solutions that could be interesting, but do not look interesting based on the chosen weights.

#### Multi-Objective Evolutionary Algorithms

• Consider different objectives separately, instead of combining them into a single fitness function.

• Advantages:

- No need to set weights.
- Having objectives in different scales is not a problem.
- Retrieve a set of solutions with different trade-offs. E.g.:



#### Multi-Objective Evolutionary Algorithms

• Consider different objectives separately, instead of combining them into a single fitness function.

• Advantages:

- No need to set weights.
- Having objectives in different scales is not a problem.
- Retrieve a set of solutions with different trade-offs.
- Disadvantage:
  - It is difficult to visualise trade-offs among >3 objectives.
  - Some multi-objective evolutionary algorithms also struggle to produce good solutions when we have >3 objectives.

#### Non-dominated Sorting Genetic Algorithm II (NSGA-II)

- One of the most famous multi-objective optimisation algorithms.
- Evolutionary algorithms put some selective pressure towards better individuals (parents and / or survival selection).
- The concept of what is a better individual is simple when there is a single objective.
  - Minimisation problem:

```
f(sol_A) = 5,

f(sol_B) = 10 \longrightarrow sol_A is better than sol_B

f(sol_A) = 5,

f(sol_B) = 5 \longrightarrow sol_A is equally good to sol_B
```

#### Non-dominated Sorting Genetic Algorithm II (NSGA-II)

- When more than one objective is considered separately, this idea does not work well. E.g.:
  - Minimise objective f1 and objective f2.
     f1(solA) = 10, f2(solA) = 5
     f1(solB) = 4, f2(solB) = 10
     Which solution is better?
- NSGA-II is based on the concept of dominance (and nondominance) to determine which solutions are better.

- A solution sol<sub>A</sub> dominates another solution sol<sub>B</sub> if the following conditions are satisfied:
  - sol<sub>A</sub> is equal or better than sol<sub>B</sub> in all objectives, and
  - $sol_A$  is strictly better than  $sol_B$  in at least one objective.

Candidate solution A: cost (min) = 15,000 duration (min) = 12 months Candidate solution B: cost (min) = 16,000 duration (min) = 24 months

Solution A dominates solution B.

- A solution sol<sub>A</sub> dominates another solution sol<sub>B</sub> if the following conditions are satisfied:
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Candidate solution A: cost (min) = 15,000 duration (min) = 12 months Candidate solution B: cost (min) = 16,000 duration (min) = 24 months

Solution B does not dominate solution A.

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Candidate solution A: cost (min) = 15,000 duration (min) = 12 months robustness (max) = 10 Candidate solution B: cost (min) = 15,000 duration (min) = 12 months robustness (max) = 11

Solution A does not dominate solution B.

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Solution A does not dominate solution B.

- A solution sol<sub>A</sub> dominates another solution sol<sub>B</sub> if the following conditions are satisfied:
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Candidate solution A: cost (min) = 15,000 duration (min) = 12 months robustness (max) = 10 Candidate solution B: cost (min) = 15,000 duration (min) = 12 months robustness (max) = 10

Solution B does not dominate solution A.

### **Optimal Solutions**

- Instead of having single best solutions, we have sets of best solutions with different trade-offs.
- Pareto front: set of solutions that are non-dominated by any other solution in the search space.
- NSGA-II aims at finding the Pareto front.



Consider that all possible solutions are plotted in this graph. Pareto front is shown in blue.



[Youtube video posted by DataAb: <u>https://youtu.be/</u> <u>sEEiGM9em8s</u>]

#### Non-dominated Fronts

- NSGA-II sorts solutions according to their quality based on nondominated fronts (non-dominated sorting).
- This makes it easier to compare solutions.



Consider that all solutions from a certain generation are plotted in this graph.

#### Non-dominated Fronts

Solutions in former fronts are considered better than solutions in latter fronts and will be preferred for parents / survival selection.



- Solutions in the same front do not dominate each other, and are not dominated by any solution from latter fronts.
- Some solutions in former fronts dominate solutions
   in latter fronts.

#### Non-dominated Fronts

NSGA-II considers solutions in former fronts as better than solutions in latter fronts and will prefer them during parents / survival selection.



Many solutions are non-dominated by each other.

How to decide between solutions from the same front?

#### Deciding Between Non-Dominated Solutions

• We prefer solutions in less crowded areas of a given front.



How to determine how crowded a certain solution is?

## Crowding Distance

- The crowding distance of a solution defines how crowded the region where a given solution is.
  - 1. Consider the non-dominated front where a solution is.
  - 2. Find the two solutions in either side of the solution.
  - 3. Calculate the distance between these two solutions in the objective space.



Solutions with higher crowding distance are in less crowded regions.

## Crowding Distance

- The crowding distance of a solution defines how crowded the region where a given solution is.
  - 1. Consider the non-dominated front where a solution is.
  - 2. Find the two solutions in either side of the solution.
  - 3. Calculate the distance between these two solutions in the objective space.



Crowding distance of sol<sub>A</sub>: 1-norm distance between sol<sub>B</sub> and sol<sub>C</sub>

 $|cost(sol_B) - cost(sol_C)| + |duration(sol_B) - duration(sol_C)|$ 

Crowding distance for the left-most and right-most solutions is infinite.

# Why Preferring Solutions from Less Crowded Areas?

- We aim at finding the Pareto set of optimal solutions.
- Pareto set provides several trade-offs among objectives.
- If we lack diversity, we may may miss some optimal solutions, and be unable to provide a good spread of trade-offs between different objectives.

![](_page_31_Picture_4.jpeg)

Favouring solutions from less crowded areas encourages more diversity, helping to find the Pareto set.

#### Deciding Between Solutions

- If  $sol_A$  is in a former non-dominated font than  $sol_B$ :
  - we prefer the sol<sub>A</sub>.
- If sol<sub>A</sub> and sol<sub>B</sub> are in the same non-dominated front:
  - we prefer the solution from less crowded area, i.e., with higher crowding distance.

## NSGA-II Algorithm

- 1. Set t = 0 (current generation)
- 2. Initialise population  $P_t$  with size N.
- 3. Sort  $P_t$  into different non-dominated fronts.
- 4. Determine the crowding distance of each individual in  $P_t$ .
- 5. While t < t<sub>max</sub>
  - 1. Select parents from *P<sub>t</sub>* using 2-tournament selection based on nondominated fronts and crowding distance.
  - 2. Apply crossover to generate children individuals *C* with probability *Pc*.
  - 3. Apply mutation to children individuals *C* with probability *Pm*.
  - 4.  $S < -P_t U C$ .
  - 5. Sort S in different non-dominated fronts  $F_0$  to  $F_n$ .
  - 6. Determine the crowding distance of each individual in *S*.
  - 7. Select survivors from S based on non-dominated fronts and crowding distance.
  - 8. t <--- t+1

#### NSGA-II Algorithm — Survivor Selection

Elitist survivor selection that picks the best individuals according to the non-dominated fronts and the crowding distance.

- 1. Set *P*<sub>*t*+1</sub> <--- {}, i <--- 0
- 2. While size of  $P_{t+1}$  + size of  $F_i \le N$

1.  $P_{t+1} < -- P_{t+1} \cup F_i$ 

2. i++

3. Top  $P_{t+1}$  up with the individuals from  $F_i$  that have the highest crowding distance.

While whole front  $F_i$  fits within  $P_{t+1}$ .

### Further Reading

A fast and elitist multiobjective genetic algorithm: NSGA-II

K. Deb, A. Pratap, S. Agarwal, T. Meyarivan

IEEE Transactions on Evolutionary Computation

Vol 6, Issue 2, pages 182-197, 2002

Read until section III.

http://ieeexplore.ieee.org/xpls/abs\_all.jsp?arnumber=996017&tag=1

#### Optional:

Multi-Objective Approaches to Optimal Testing Resource Allocation in Modular Software Systems

Z. Wang, K. Tang, X. Yao

IEEE Transactions on Reliability

Vol 59, Issue 3, pages 563—575, 2010

http://ieeexplore.ieee.org/abstract/document/5549979/